Performance Analysis of Coarse Solvers for Algebraic Multigrid on Leading Multicore Architectures

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Outline

- New spectral element AMG method
 - Set up hierarchy
 - Coarse grid solver
- Performance characterization, multicore optimization, bound analysis
 - Intel MIC Knights Corner
- Software components







Advances in multigrid solver: Smoothed aggregation spectral element AMG

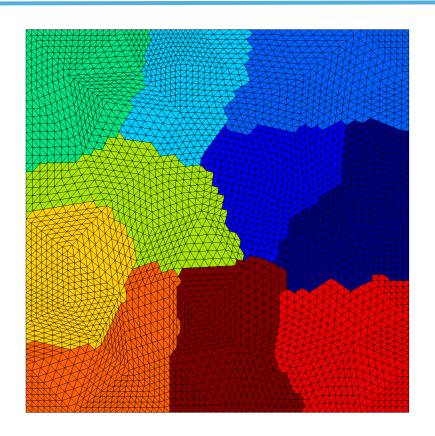
- AMGe is geometric multigrid with nonstandard elements (agglomerates of fine-grid ones) and operator-dependent coarse finite element spaces
- For elliptic problems (diffusion and elasticity), solve local eigenvalue problems to build AMG hierarchy – highly parallel
 - Code SAAMGe released
- For more general problems (electromagnetics and Darcy flow), solve local SVD problems to build AMG hierarchy – highly parallel
 - Prototype code ParElag
- Novelty: The hierarchy can be employed for nonlinear solvers on unstructured meshes as well as for MCMC simulations
 - In contrast to plain AMG, AMGe coarse spaces have guaranteed approximation properties so the coarse problems provide highly accurate discretizations (useful for nonlinear problems and for dimension reduction)



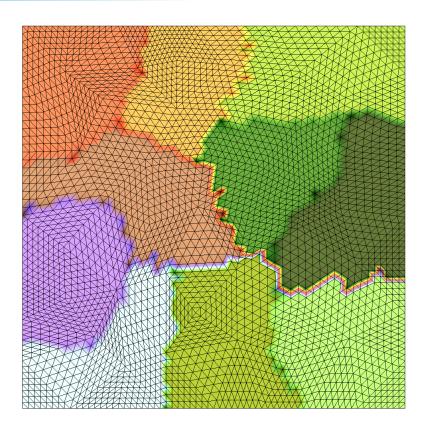




Agglomeration and Aggregation in finite elements



elements → agglomerates



vertices → aggregates







Smoothed aggregation spectral element AMG (SAAMGE)

- High order elements to discretize PDEs
- Algebraic multigrid (AMG) solver
 - 1. Pre-smoothing (fine grid)
 - Intermediate iterate: $y = x_i + M^{-1}(b A x_i)$
 - 2. Coarse-grid correction (recursion → multilevel)
 - 1) Restrict the residual: $r_c = P^T (b A y)$

PCG, hypre,
HSS, . . .

- 2) Solve coarse-grid defect equation: $A_c x_c = r_c$
- 3) Interpolate, compute next intermediate iterate: $z = y + P x_c$
- 3. Post-smoothing (fine grid)
 - $x_{i+1} = z + M^{-T} (b A z)$







Interpolation matrix P for coarse space construction (set up AMG hierarchy)

Coarse-grid matrix $A_c = P^TAP$; P is computed from lowest eigenvectors of local element matrices

 Improves approximation quality of coarse space, but expensive to compute

$$P = \begin{vmatrix} P_1 \\ P_2 \\ \vdots \\ P_{na} \end{vmatrix}$$

Tentative $P = \begin{bmatrix} P_1 \\ P_2 \\ P_{na} \end{bmatrix}$ P_i is obtained by solving eigenvalue problem A_i q = λ D_i q, need several smallest eigenpairs

Final P = S P, where S is a matrix polynomial (e.g. Chebyshev)

$$P \qquad S \qquad \hat{P}$$

$$A_c \qquad P^T \qquad A \qquad P$$







Interpolation matrix P for coarse space construction

Two improvements:

- Replace dense LAPACK by sparse ARPACK + SuperLU
 - Projection method reduces to smaller problems
- Developed an early termination scheme by monitoring the accuracy of already computed eigenpairs
 - Using implicit QL method for tridiagonal matrices
- → over 2x faster in total solution time

Osni Marques, Wednesday, 10:50 am,

"Tuning the Coarse Space Construction in a Spectral AMG Solver"







Advances in low-rank HSS factorization solver

- Goal: achieve O(N), or O(N polylog(N)) complexity
 - Traditional sparse LU requires O(N²) Flops (e.g., SuperLU)
- Approach: use hierarchical matrix algebra
 - Accurate approximation with low-rank, data-sparse structures
- Same mathematical foundation as Fast Multipole (FMM), but more general
 - Diagonal block ("near field") exact; off-diagonal block ("far field") approximated via low-rank format

$$A \approx \left[\begin{array}{cc} D_1 & U_1 B_1 V_2^T \\ U_2 B_2 V_1^T & D_2 \end{array} \right]$$

- Broad applications
 - PDEs with smooth kernels, Integral equations, Boundary Element Methods, genreal-purpose preconditioners, ...

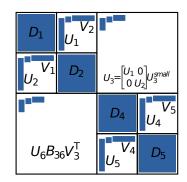




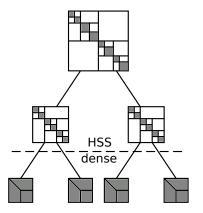


Use hierarchical partitioning and nested bases to achieve lower complexity

- STRUMPACK, ~200 downloads in 2015
 - 5.4x faster than dense LU for BEM matrices
 - 7x faster than traditional sparse solver for PDEs, 4-fold memory reduction
 - CEMM fusion SciDAC problems (two-fluid MHD)
 - ComPASS accelerator SciDAC problems (Maxwell equations)



<u>Pieter Ghysels, Wednesday, 5:45 pm</u>
"Evaluation of a preconditioner using low-rank approximation and randomized sampling"





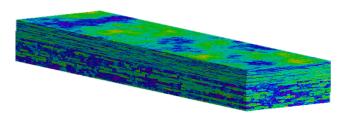




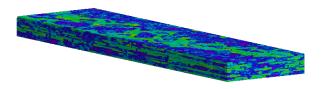
Benchmark: PDE with jump coefficients

SPE10 Benchmark (model 2) (www.spe.org/web/csp/)

- Formation in the Brent oil field; 1200'x2200'x170' (cell size 20'x10'x2').
- Top 70 ft (35 layers) represents the Tarbert formation; bottom 100 ft (50 layers) represents Upper Ness (fluvial).



porosity of the whole model

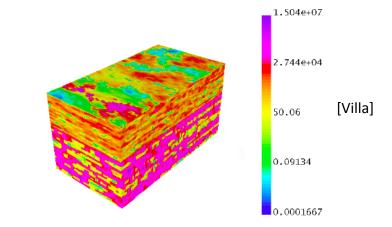


part of the Upper Ness sequence

Darcy equation

$$-\nabla \cdot (k(x)\nabla p) = f(x), \forall x \in \Omega$$

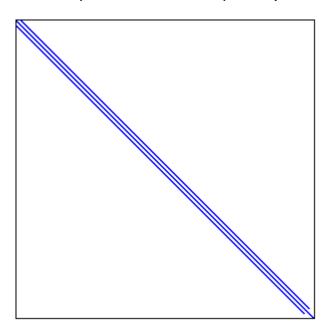
 $p(x)$ is pressure
 $k(x)$ is permeability of the medium



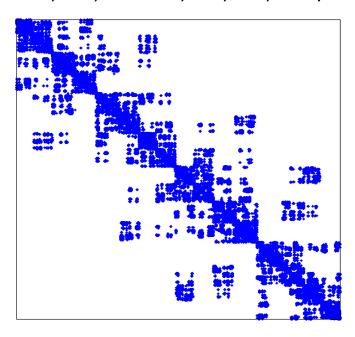
Two distinct soil layers and the large jumps in the coefficient k between them.



Fine-grid, coarse-grid matrices from SAAMGE



 $n \approx 1.2 \text{ M}$; $nnz \approx 30.6 \text{ M}$; $nnz/n \approx 26$ n = 7,782; nnz = 1,412,840; nnz/n = 181.6



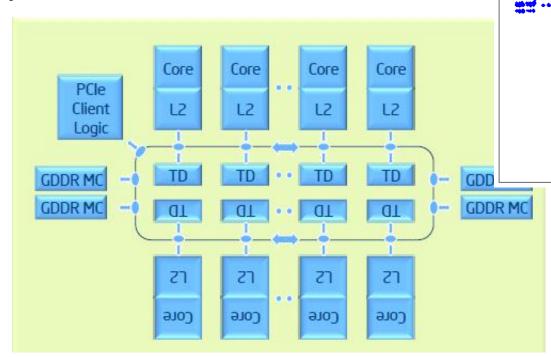




Multicore example

Intel Manycore Integrated Core (MIC), Knights Corner

- 60 cores per card
- 4 hardware threads / core
- 512-bit SIMD Vector = 8 DP FLOPS / cycle
- L2 private cache, coherence via bidirection.

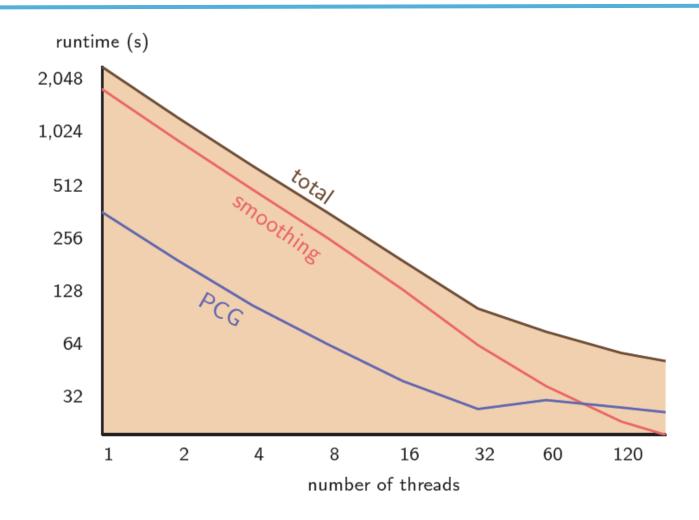








Time breakdown of the AMG cycle









PCG thread-friendly optimization

Goal: minimize threads synchronization

Algorithm 1

1: while not converged do

2: $\rho \leftarrow \sigma$

3: **omp parallel for :** $w \leftarrow Ap$

4: **omp parallel for :** $\tau \leftarrow w \cdot p$

5: $\alpha \leftarrow \rho / \tau$

6: **omp parallel for :** $x \leftarrow x + \alpha p$

7: **omp parallel for :** $r \leftarrow r - \alpha w$

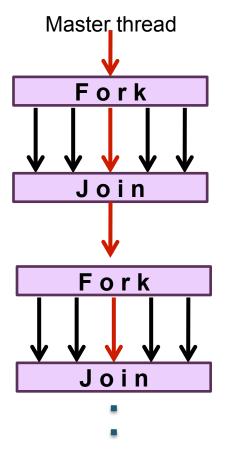
8: **omp parallel for :** $z \leftarrow M^{-1}r$

9: **omp parallel for :** $\sigma \leftarrow z \cdot r$

10: $\beta \leftarrow \sigma / \rho$

11: **omp parallel for :** $p \leftarrow z + \beta p$

12: end while









Improve PCG: Algorithm 2 (omp-for-all)

1: omp parallal

 \mapsto single parallel region

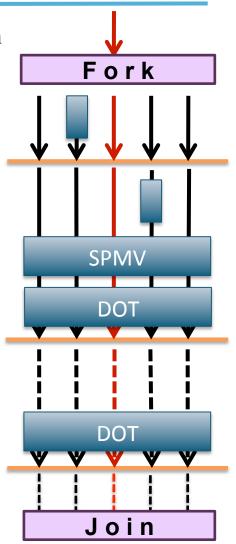
- while not converged do
- omp single: $\tau \leftarrow 0.0$ 3:

 \mapsto implied barrier

- omp single nowait : $\rho \leftarrow \sigma, \sigma \leftarrow 0.0$ 4:
- 5: omp for nowait: $w \leftarrow A p$
- **omp for reduction :** $\tau \leftarrow w \cdot p \mapsto \text{implied barrier}$ 6:

- 7: $\alpha \leftarrow \rho / \tau$
- 8: **omp for nowait :** $x \leftarrow x + \alpha p$
- 9: **omp for nowait :** $r \leftarrow r - \alpha w$
- **omp for nowait :** $z \leftarrow M^{-1}r$ 10:
- 11:
 - **omp for reduction:** $\sigma \leftarrow z \cdot r \mapsto \text{implied barrier}$
- 12: $\beta \leftarrow \sigma / \rho$
- **omp for nowait :** $p \leftarrow z + \beta p$ 13:
- 14: end while
- 15: end omp parallel







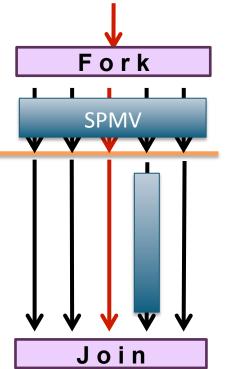


Improve PCG: Algorithm 3 (omp-for-spmv)

1: omp parallal

 \mapsto single parallel region

- 2: **while** not converged **do**
- 3: **omp for :** $w \leftarrow Ap$
- 4: **omp single**
- 5: $\tau \leftarrow w \cdot p$
- 6: $\alpha \leftarrow \rho / \tau$
- 7: $x \leftarrow x + \alpha p$
- 8: $r \leftarrow r \alpha w$
- 9: $z \leftarrow M^{-1}r$
- 10: $\rho \leftarrow \sigma$
- 11: $\sigma \leftarrow z \cdot r$
- 12: $\beta \leftarrow \sigma / \rho$
- 13: $p \leftarrow z + \beta p$
- 14: **end omp single**
- 15: end while

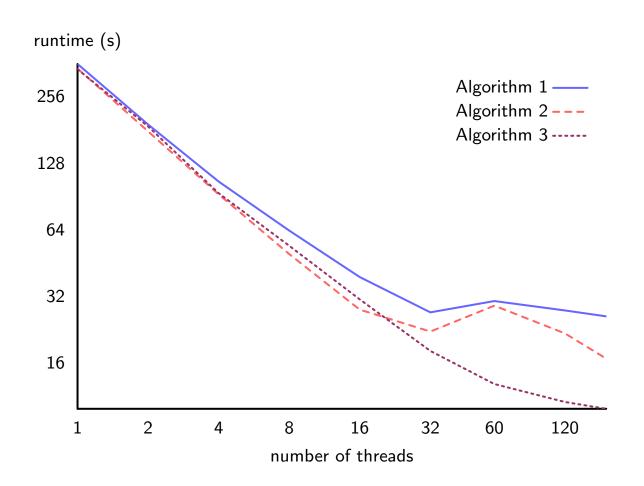








PCG improved runtime









Roofline performance modeling (collaborating with SUPER SciDAC Institute)

• Run time = MAX (memory access time, flops time)

stage	bytes	flops
pre- and post-smooth restrict coarse solve (PCG/J)	$(3\nu+1)(12\mathtt{nza}+3\cdot8n)\\12\mathtt{nza}+12\mathtt{nzp}+3\cdot8n$	$\frac{2(3\nu+1)(\mathtt{nza}+2n)}{2(\mathtt{nza}+\mathtt{nzp})}$
multiply by A_c preconditioner vector operations interpolate termination criterion	$egin{array}{l} 12\mathtt{nzc} \ 2\cdot 8n_c \ 5\cdot 8n_c \ 12\mathtt{nzp} + 8n \ 12\mathtt{nza} + 4\cdot 8n \end{array}$	$2 exttt{nzc}$ n_c $2\cdot 5n_c$ $2 exttt{nzp}$ $2(exttt{nza}+n)$

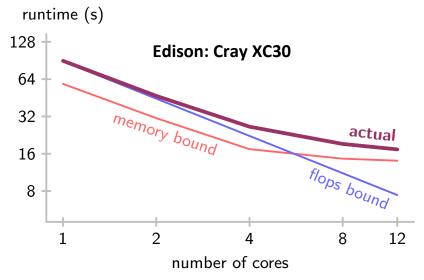




Roofline performance modeling (collaborating with SUPER SciDAC Institute)

• Run time = MAX (flops time, memory access time)

stage	bytes	flops
pre- and post-smooth restrict	$(3\nu+1)(12{\tt nza}+3\cdot 8n) \ 12{\tt nza}+12{\tt nzp}+3\cdot 8n$	$\frac{2(3\nu+1)(\mathtt{nza}+2n)}{2(\mathtt{nza}+\mathtt{nzp})}$
coarse solve (PCG/J) multiply by A_c preconditioner	$12 exttt{nzc}$ $2\cdot 8n_c$	$2\mathtt{nzc}$ n_c
vector operations interpolate termination criterion	$5 \cdot 8n_c$ $12 \operatorname{nzp} + 8n$ $12 \operatorname{nza} + 4 \cdot 8n$	$2 \cdot 5n_c$ $2 \mathtt{nzp}$ $2(\mathtt{nza} + n)$



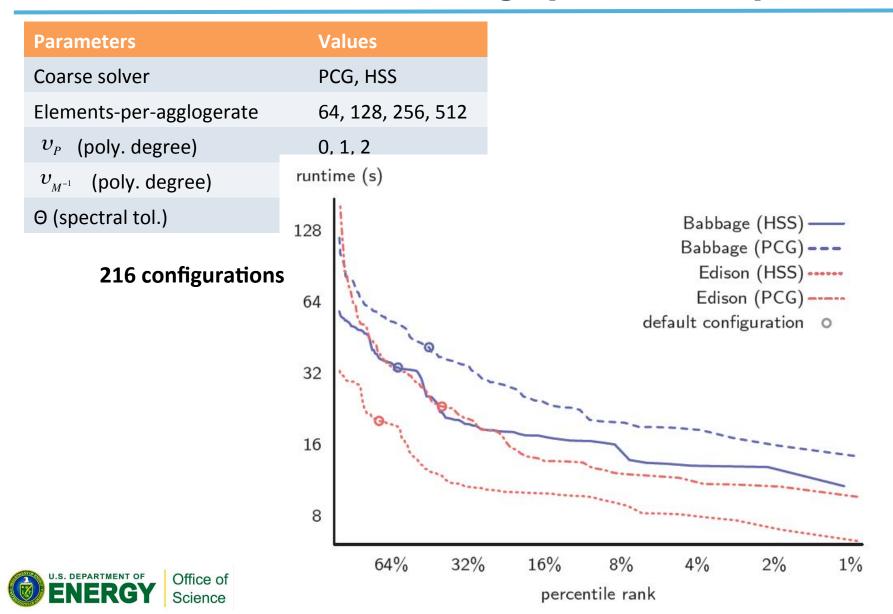
Science

Roofline bound gap

Edison: 12- cores	Coarse PCG	Coarse HSS
Gap from bound	23%	31%

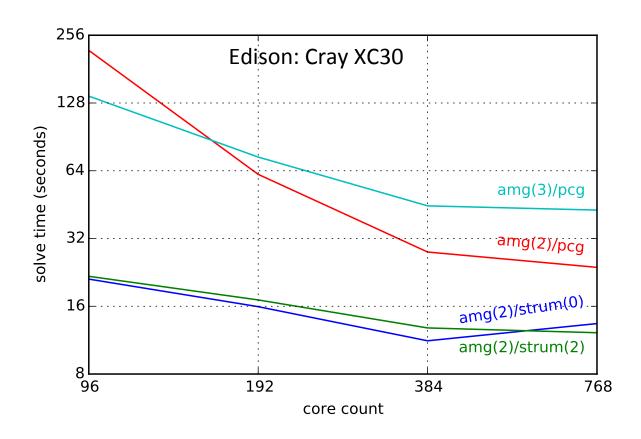
A. Druinsky, P. Ghysels, X.S. Li, O. Marques, S. Williams, A. Barker, D. Kalchev, P. Vassilevski,. "Comparative Performance Analysis of Coarse Solvers for Algebraic Multigrid on Leading Multicore Architectures", PPAM 2015.

Performance variation in large paramater space



Parallel performance of multilevel AMGe + HSS

- SPE10, 8 million DOFs
- Speed up MG time up to 6.4x after HSS replaced PCG as coarse level solver



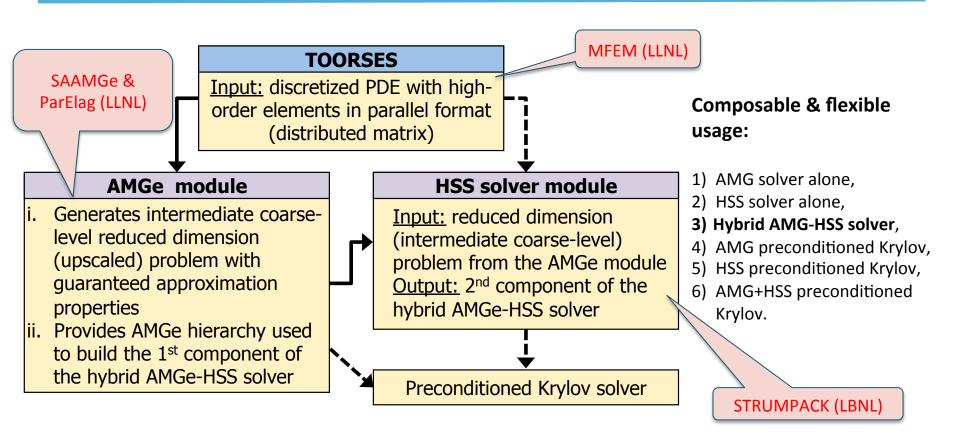






TOORSES software stack

(Towards Optimal Order Resilient Solvers at Extreme Scale)



- MFEM: http://mfem.org
- SAAMGe, ParElag: https://myconfluence.llnl.gov/display/AMGE/AMGe
- STRUMPACK: http://portal.nersc.gov/project/sparse/strumpack







Summary, on-going work

- Fix scalability issues in the multilevel MPI code
 - Comprehensive performance evaluation
- Solvers' parameters greatly influence performance
 - Develop autotuning framework to help parameter selection
- Initial success of two-grid, single node roofline model to understand performance limit
 - Extend to multilevel code, multiple nodes with MPI communication
 - Improve roofline model for factorization algorithms





